

CS 161

Intro to CS I

Finish Recursion/Begin Memory Model



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Odds and Ends

- Assignment 5 posted
- Assignment 4 demo this week
- 1 credit hour of lecture/course is 3 hours outside the course
- Poor planning on your part does not constitute an emergency on mine!
- KISS!

Iterative Factorial



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```
factorial(0) = 1;  
factorial(n) = n*n-1*n-2*...*n-(n-1)*1;
```

```
long factorial(int n) {  
    long fact;  
    if(n==0)  
        fact=1;  
    else  
        for(fact=n; n > 1; n--)  
            fact=fact*(n-1);  
    return fact;  
}
```

Recursive Factorial



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factorial(0) = 1; *base case*
factorial(n) = n * factorial(n-1);

```
long factorial(int n) {  
    if (n == 0)      // Base case  
        return 1;  
    else  
        return n * factorial(n - 1); // Recursive call  
}
```

Computing Factorial Iteratively



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factorial(4)

$$\text{factorial}(0) = 1;$$

$$\text{factorial}(n) = n * (n-1) * \dots * 2 * 1;$$

Computing Factorial Iteratively



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$\text{factorial}(4) = 4 * 3$

```
factorial(0) = 1;  
factorial(n) = n*(n-1)*...*2*1;
```

Computing Factorial Iteratively



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$$\text{factorial}(4) = \underbrace{4 * 3}_{= 12} * 2$$

```
factorial(0) = 1;  
factorial(n) = n*(n-1)*...*2*1;
```

Computing Factorial Iteratively



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$$\begin{aligned}\text{factorial}(4) &= 4 * 3 \\ &= \underline{12} * 2 \\ &= \cancel{24} * 1\end{aligned}$$

```
factorial(0) = 1;  
factorial(n) = n*(n-1)*...*2*1;
```

Computing Factorial Iteratively



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$$\begin{aligned}\text{factorial}(4) &= 4 * 3 \\&= 12 * 2 \\&= 24 * 1 \\&= 24\end{aligned}$$

```
factorial(0) = 1;  
factorial(n) = n*(n-1)*...*2*1;
```

Computing Factorial Recursively



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factorial(4)

$$\text{factorial}(0) = 1;$$

$$\text{factorial}(n) = n * \text{factorial}(n-1);$$

Computing Factorial Recursively



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```
factorial(0) = 1;  
factorial(n) = n * factorial(n-1);
```

$\text{factorial}(4) = 4 * \underline{\text{factorial}(3)}$

Computing Factorial Recursively



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```
factorial(0) = 1;  
factorial(n) = n * factorial(n-1);
```

$$\begin{aligned}\text{factorial}(4) &= 4 * \text{factorial}(3) \\ &= \underline{4} * (3 * \text{factorial}(2))\end{aligned}$$

Computing Factorial Recursively



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```
factorial(0) = 1;  
factorial(n) = n*factorial(n-1);
```

$$\begin{aligned}\text{factorial}(4) &= 4 * \text{factorial}(3) \\ &= 4 * (3 * \text{factorial}(2)) \\ &= 4 * (3 * (2 * \text{factorial}(1)))\end{aligned}$$

Computing Factorial Recursively



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```
factorial(0) = 1;  
factorial(n) = n * factorial(n-1);
```

$$\begin{aligned}\text{factorial}(4) &= 4 * \text{factorial}(3) \\&= 4 * (3 * \text{factorial}(2)) \\&= 4 * (3 * (2 * \text{factorial}(1))) \\&= 4 * (3 * (2 * (1 * \text{factorial}(0))))\end{aligned}$$

Computing Factorial Recursively



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```
factorial(0) = 1;  
factorial(n) = n * factorial(n-1);
```

$$\begin{aligned}\text{factorial}(4) &= 4 * \text{factorial}(3) \\&= 4 * (3 * \text{factorial}(2)) \\&= 4 * (3 * (2 * \text{factorial}(1))) \\&= 4 * (3 * (2 * (1 * \text{factorial}(0)))) \\&= 4 * (3 * (2 * (1 * 1)))\end{aligned}$$

↙

Computing Factorial Recursively



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```
factorial(0) = 1;  
factorial(n) = n*factorial(n-1);
```

$$\begin{aligned}\text{factorial}(4) &= 4 * \text{factorial}(3) \\&= 4 * (3 * \text{factorial}(2)) \\&= 4 * (3 * (2 * \text{factorial}(1))) \\&= 4 * (3 * (2 * (1 * \text{factorial}(0)))) \\&= 4 * (3 * (2 * (1 * 1))) \\&= 4 * (3 * (2 * 1))\end{aligned}$$

Computing Factorial Recursively



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```
factorial(0) = 1;  
factorial(n) = n * factorial(n-1);
```

$$\begin{aligned}\text{factorial}(4) &= 4 * \text{factorial}(3) \\&= 4 * (3 * \text{factorial}(2)) \\&= 4 * (3 * (2 * \text{factorial}(1))) \\&= 4 * (3 * (2 * (1 * \text{factorial}(0)))) \\&= 4 * (3 * (2 * (1 * 1))) \\&= 4 * (3 * (2 * 1)) \\&= 4 * (3 * 2)\end{aligned}$$

Computing Factorial Recursively



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```
factorial(0) = 1;  
factorial(n) = n * factorial(n-1);
```

$$\begin{aligned}\text{factorial}(4) &= 4 * \text{factorial}(3) \\&= 4 * (3 * \text{factorial}(2)) \\&= 4 * (3 * (2 * \text{factorial}(1))) \\&= 4 * (3 * (2 * (1 * \text{factorial}(0)))) \\&= 4 * (3 * (2 * (1 * 1))) \\&= 4 * (3 * (2 * 1)) \\&= 4 * (3 * 2) \\&= 4 * 6\end{aligned}$$

Computing Factorial Recursively



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$$\begin{aligned}\text{factorial}(4) &= 4 * \text{factorial}(3) & \text{factorial}(0) &= 1; \\ &= 4 * (3 * \text{factorial}(2)) & \text{factorial}(n) &= n * \text{factorial}(n-1); \\ &= 4 * (3 * (2 * \text{factorial}(1))) \\ &= 4 * (3 * (2 * (1 * \text{factorial}(0)))) \\ &= 4 * (3 * (2 * (1 * 1))) \\ &= 4 * (3 * (2 * 1)) \\ &= 4 * (3 * 2) \\ &= 4 * 6 \\ &= 24\end{aligned}$$

Differences



- Pros
 - Readability
- Cons
 - Efficiency
 - Memory

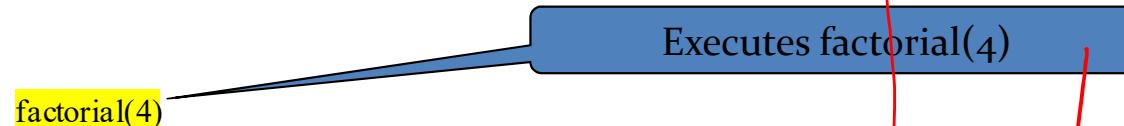
You have to have shallow depth
but too many recursive calls.

Recursive Factorial



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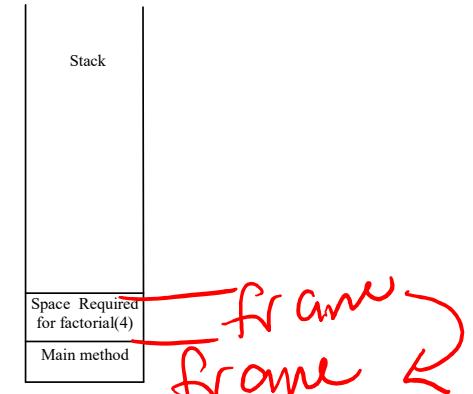
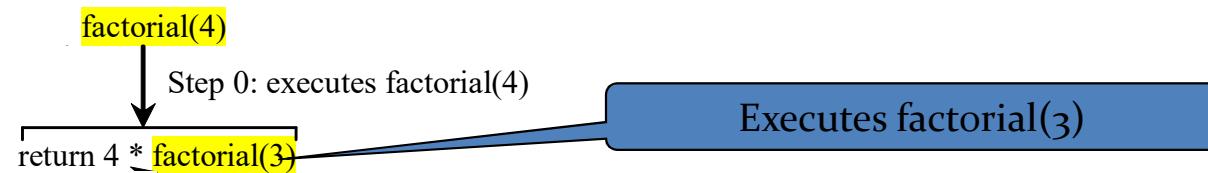
recap



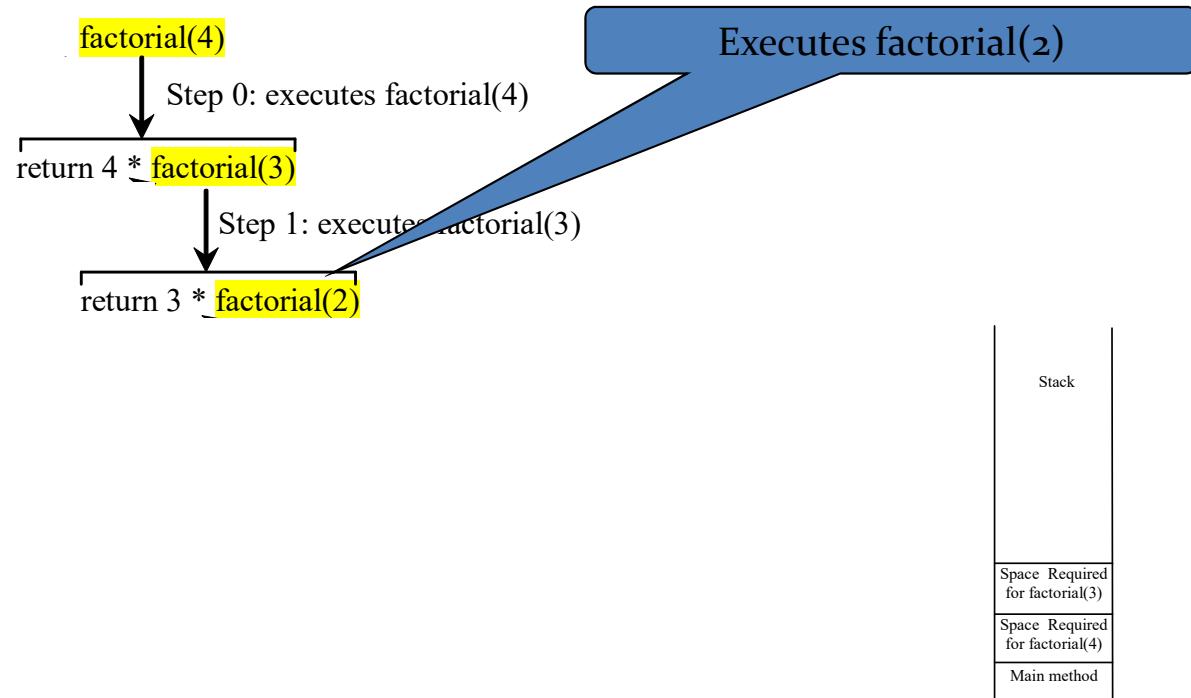


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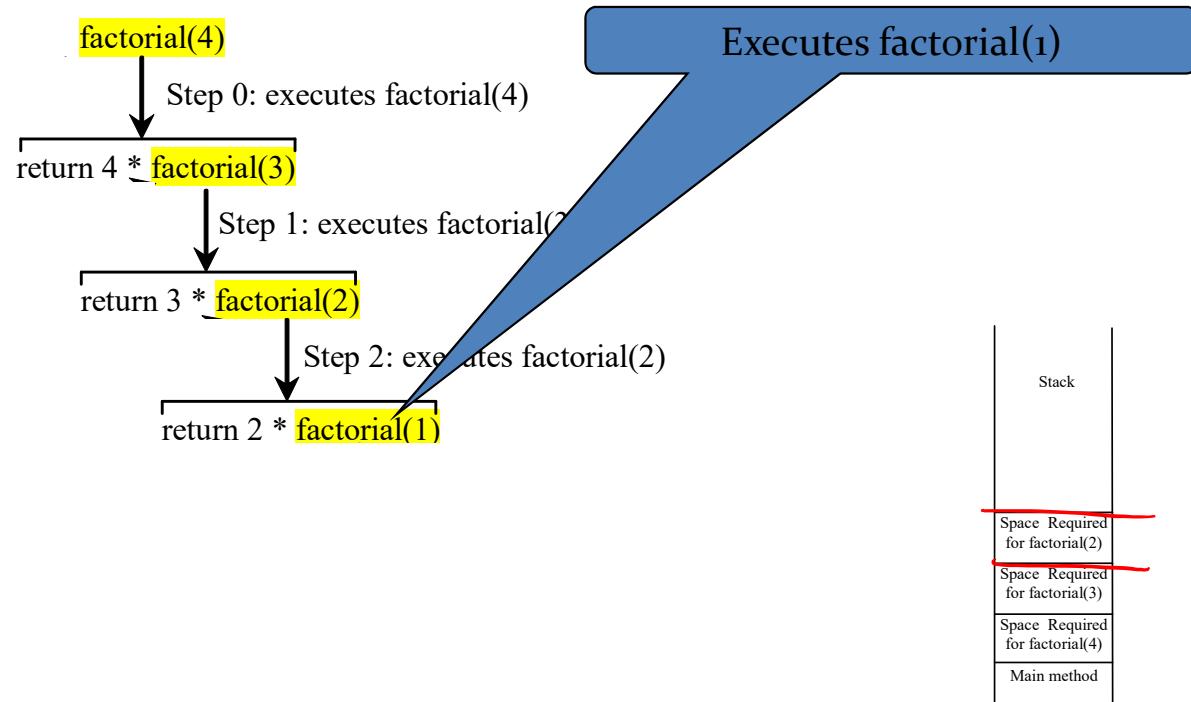
Recursive Factorial



Recursive Factorial



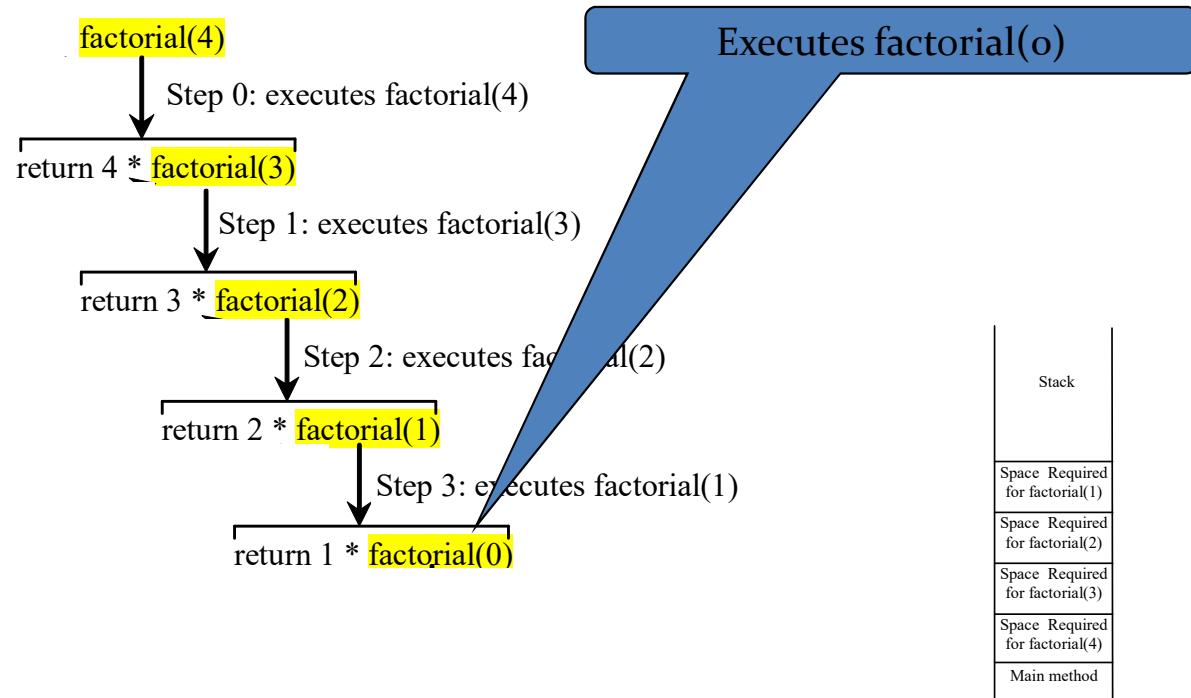
Recursive Factorial



Recursive Factorial



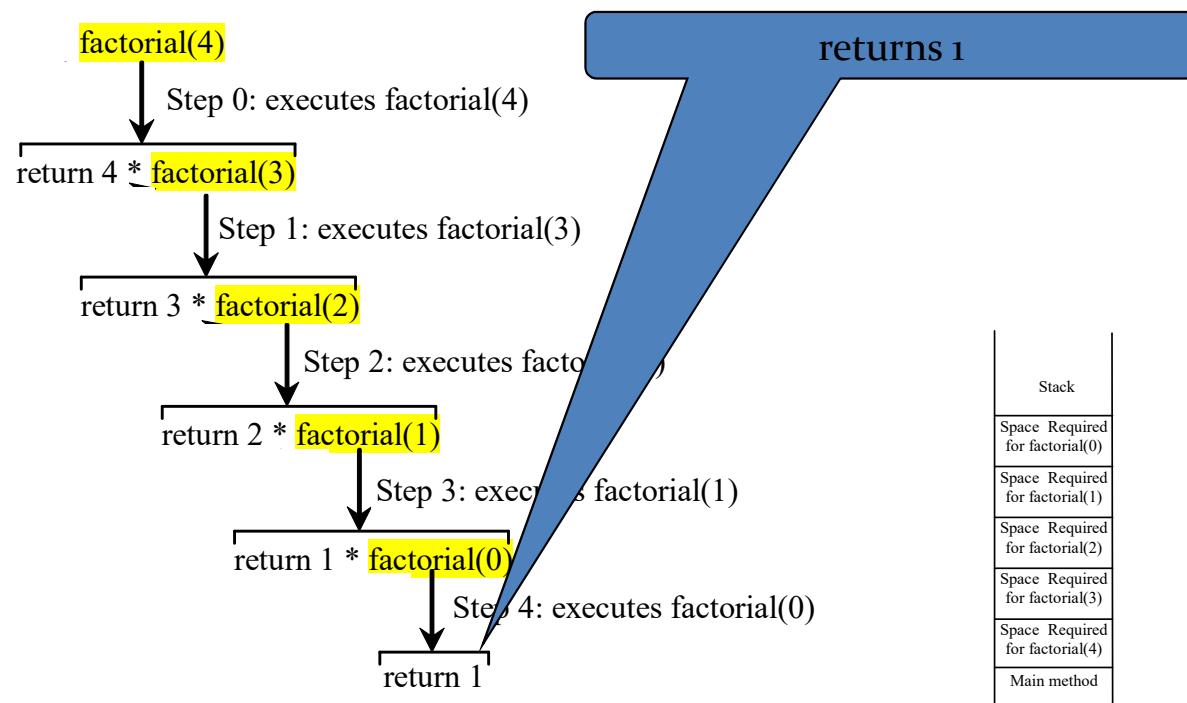
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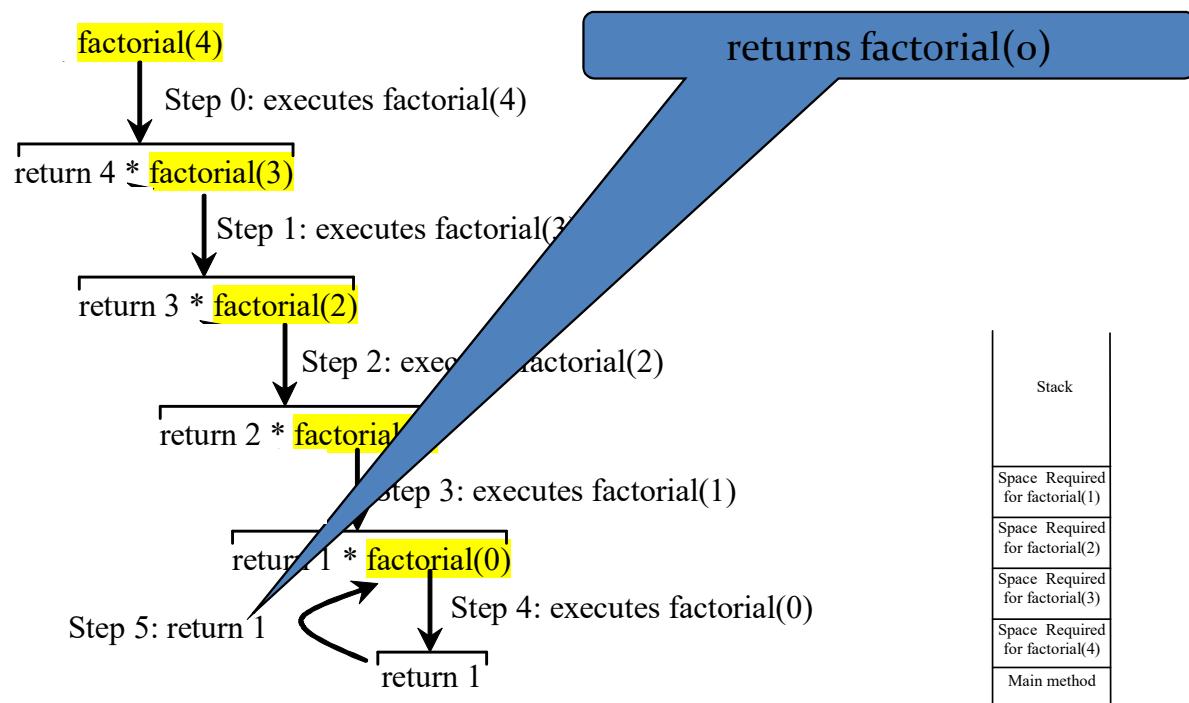
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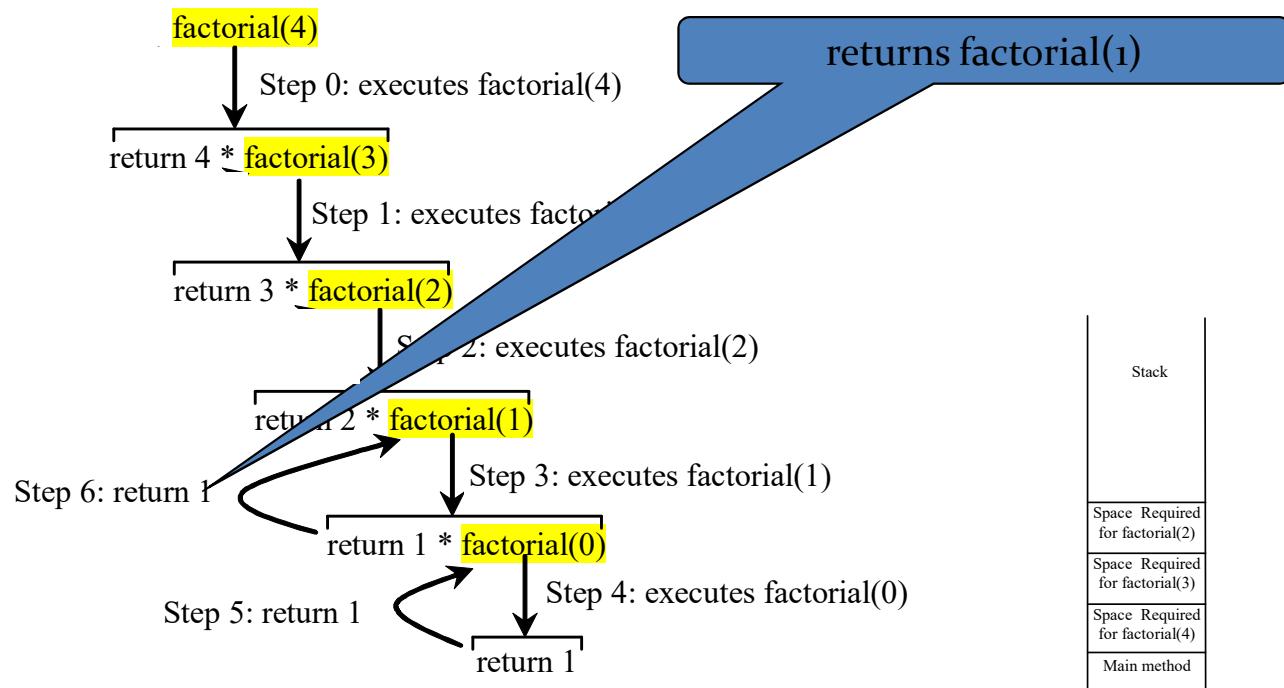
Recursive Factorial



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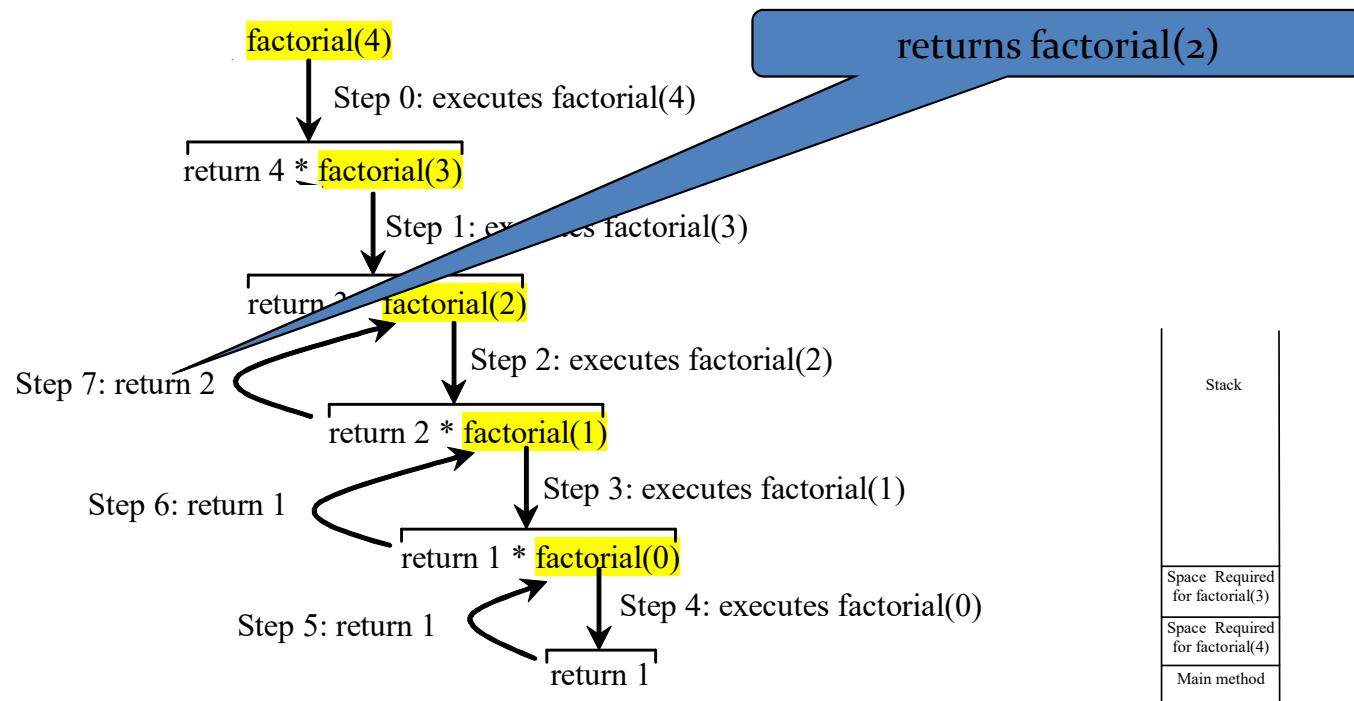
Recursive Factorial



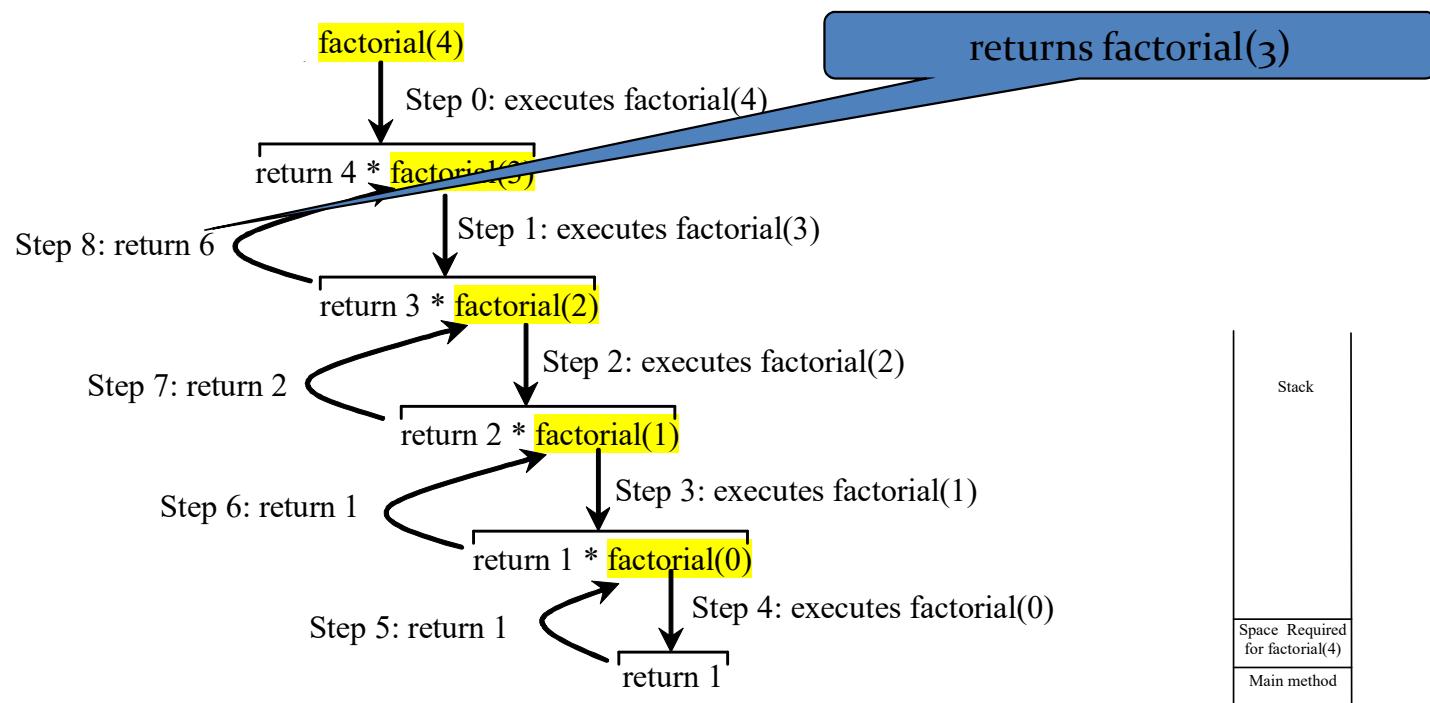
Recursive Factorial



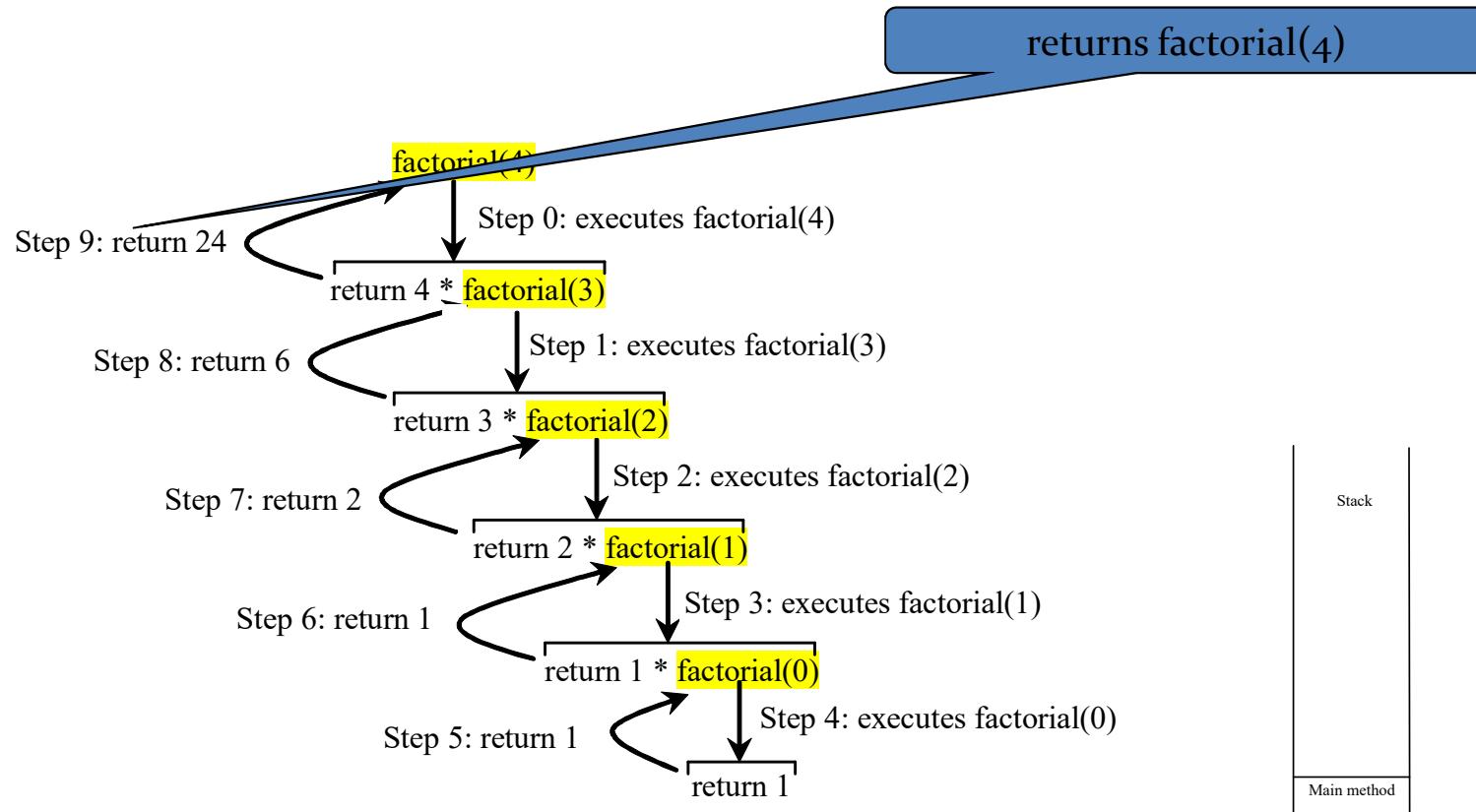
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Recursive Factorial



Recursive Factorial



In-class Exercise #4



- ~~Get into groups of 4 - 5.~~
- Write your own recursive *int pwr()* function that takes two integers as arguments and returns the integer result.
 - What does the function prototype look like?
 - Now, write the function definition...



4. ENGR

Re-attach Fullscreen Stay on top Duplicate Close

```
1 #include <iostream>
2
3 using namespace std;
4
5 int pwr(int base, int exp) {
6     int result=1;
7
8     for(int i=0; i<exp; i++) {
9         result=result*base;
10    }
11
12    return result;
13 }
14 int pwr_r(int base, int exp) {
15     if(exp==0) //base case
16         return 1;
17     else
18         return base*pwr_r(base, exp-1); //get us closer to base case
19 }
20
21 int main() {
22
23     cout << pwr(2,1000000) << endl;
24     cout << pwr_r(2,1000000) << endl; //too deep of recursion, blow stack
25
-- INSERT --
```

24,74

Top

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Stack vs. Heap

- Static vs. Dynamic

created/memory is known about
at compilation

int i;
int x;

created at
runtime

heap

You
can

You
can't
control

stack

nameless

initial
item
created
has a
name

Static vs. Dynamic

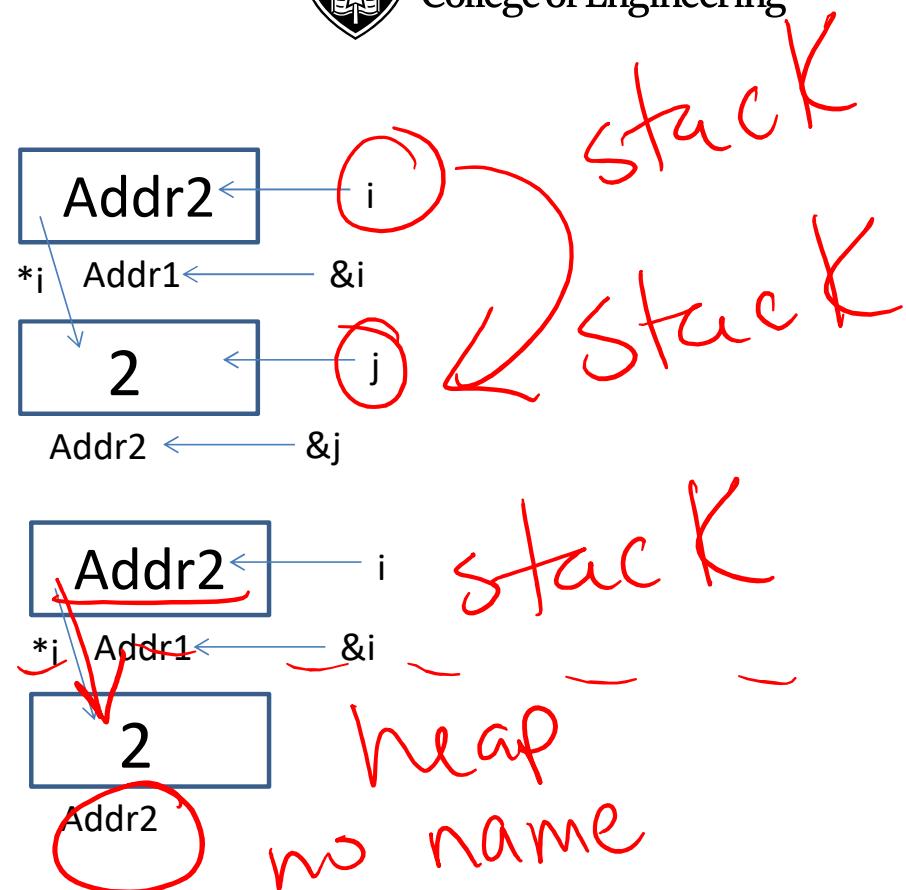
- Static Semantics
 - Assign address of variable

```
int *i, j=2;  
i=&j;
```

- Dynamic Semantics
 - Create memory
 - Assign memory to pointer

```
int *i=NULL;  
i=new int;  
*i=2;
```

extra addrs
what type
make mem on the heap





What About Memory Leaks?

- What happens here...

...

```
int main () {  
    int *i=NULL;    //created in main function  
    while(1) {  
        i = new int;  
    }  
}
```

Fixing Memory Leaks...



- What happens here...

...

```
int main () {  
    int *i=NULL;    //created in main function  
    while(1) {  
        i = new int;  
        delete i; //free memory that i points to, preventing mem leaks  
    }  
}
```

cout << i

The code shows a memory leak. The variable *i* is dynamically allocated within a loop, but the corresponding deallocation statement `delete i;` is missing. A red annotation *cout << i* is written above the loop, with arrows pointing from the *cout* to the *i* in the allocation line and from the *i* in the allocation line to the *i* in the deallocation line. The entire deallocation line is crossed out with a large red line.