CS 444/544 OS II Lab Tutorial #2 (part 1)

Booting and GDB Practice

Acknowledgement: Slides drawn heavily from Yeongjin Jiang

Lab Setup Check

Output/messages	s				
[0:7c00] => 0x7					
Breakpoint 1, 0x00	007c00 in ?? ()				
— Registers ——					
eax 0x0000aa55	ecx 0x00000000	edx 0x00000080	ebx 0x00000000		
esp 0x00006f20	ebp 0x00000000	esi 0x0000000	edi 0x00000000		
eip 0x00007c00	eflags [IF]	cs 0x00000000	ss 0x00000000		
ds 0x00000000	es 0x00000000	fs 0x00000000	gs 0x00000000		
Assembly					
0x00007c00 ? cli					
0x00007c01 ? cld					
0x00007c02 ? xor	%ax,%ax				
0x00007c04 ? mov	%ax,%ds				
0x00007c06 ? mov	%ax,%es				
0x00007c08 ? mov	%ax,%ss				
0x00007c0a ? in	\$0x64, <mark>%al</mark>				
Source					
Stack					
[0] from 0x00007c0	0				
(no arguments)					
Memory					
— Expressions —					
>>>					

Contents

- Following the Booting Sequence
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- Coding Style

Exercise 3

Set a breakpoint at address 0x7c00, which is where the boot sector will be loaded. Continue execution until that breakpoint. Trace through the code in **boot/boot.5**, using the source code and the disassembly file <a>obj/boot/boot.asm to keep track of where you are. Also use the x/i command in GDB to disassemble sequences of instructions in the boot loader, and compare the original boot loader source code with both the disassembly in <a>obj/boot/boot.asm and GDB.

Exercise 3: How?

• Use tmux, open boot/boot.S and gdb at the same time..

Use ni,si, breakpoint

bas	sh • \$1 ×	bash 362 × ssh 🔹 363	× ssh 364				
b/boot.	.S		buffers	Output/message	es		
8 .set	PROT_MODE_CSE	G, 0x8 # kernel code se	gment selector	0:7c00] => 0x7	7c00: cli		
9 .set	PROT_MODE_DSE						
1000	t CR0_PE_ON,	0x1 # protected mode	enable flag	reakpoint 1, 0x00	0007c00 in ?? ()		
11				— Registers —			
	obl start			eax 0x0000aa55		edx 0x0000080	ebx 0x0000000
13 star				esp 0x00006f20		esi 0x00000000	edi 0x00000000
and the second se	code16	# Assemble for 16-b		eip 0x00007c00		cs 0x0000000	ss 0x00000000
15 cl		# Disable interrupt		ds 0x00000000	es 0x0000000	fs 0x00000000	gs 0x00000000
16 cl	La	# String operations		Assembly			
17 18 #	Cot un the imm	ortant data segment registers (x00007c01 ? cld			
a second s	orw %ax,%ax	# Segment number ze		x00007c02 ? xor	%ax,%ax		
	ovw %ax,%ds	# -> Data Segment		x00007c04 ? mov	%ax,%ds		
-	ovw %ax,%es	# -> Extra Segment		x00007c06 ? mov	%ax,%es		
10.000	ovw %ax,%ss	# -> Stack Segment		x00007c08 ? mov	%ax,%ss		
23	orth Marcy 1000			x00007c0a ? in	\$0x64,%al		
	Enable A20:			- Source			
25 #	For backward	s compatibility with the earlie	st PCs, physical	Stack			
26 #	address line	20 is tied low, so that addres	ses higher than	0] from 0x00007c	00		
27 #	1MB wrap aro	und to zero by default. This c	ode undoes this.	no arguments)			
28 seta	a20.1:			- Memory			
29 ir	nb \$0x64,%a	l # Wait for not		 Expressions — 			
1000	estb \$0x2,%al						
31 jr 32	nz seta20.1		;	·>>			
33 mc	ovb \$0xd1,%a		0x64				
34 o l 35	utb %al,\$0x6	4					
36 seta	20.2:						
37 ir		1 # Wait for not	busy				
	estb \$0x2,%al						
39 jr	nz seta20.2						
40 I boo	ot/boot.S	asn	20% ≡ 17: 1				
- INSER							
: vim	0:vim* 1:make-					[04,	/09/2019 07:01AM]

Exercise 3: Enabling Protected Mode

call bootmain

PROT_MODE_CSEG 0x8

Bootmain() is in boot/main.c

lgdt movl orl movl	\$CR0_PE_ON, %eax	
	to next instruction, ches processor into 32	but in 32-bit code segment. 2-bit mode.
ljmp	<pre>\$PROT_MODE_CSEG, \$pr</pre>	otcseg
.code32	2	<pre># Assemble for 32-bit mode</pre>
protcseg	•	
# Set u	up the protected-mode	data segment registers
mo∨w	<pre>\$PROT_MODE_DSEG, %ax</pre>	# Our data segment selector
mo∨w	%ax, %ds	# -> DS: Data Segment
mo∨w	%ax, %es	# -> ES: Extra Segment
mo∨w	%ax, %fs	# -> FS
mo∨w	%ax, %gs	# -> GS
mo∨w	%ax, %ss	<pre># -> SS: Stack Segment</pre>
# Set ι	up the stack pointer a	nd call into C.
movl	\$start, %esp	
call be	notmain	

Exercise 3: bootmain

	_		
0x00007c45	?	call	0x7d0a
0x00007c4a	?	jmp	0x7c4a
0x00007c4c	?	add	%al,(%eax)
0x00007c4e	?	add	%al,(%eax)
0x00007c50	?	add	%al,(%eax)
0x00007c52	?	add	%al,(%eax)
0x00007c54	?	(bad)	

In boot/main.c

Set up the stack pointer and call into C. \$start, %esp mo∨l call bootmain void bootmain(void) struct Proghdr *ph, *eph; // read 1st page off disk readseg((uint32_t) ELFHDR, SECTSIZE*8, 0); // is this a valid ELF? if (ELFHDR->e_magic != ELF_MAGIC) goto bad; // load each program segment (ignores ph flags) ph = (struct Proghdr *) ((uint8_t *) ELFHDR + ELFHDR->e_phoff); eph = ph + ELFHDR->e_phnum; for (; ph < eph; ph++) // p_pa is the load address of this segment (as well // as the physical address) readseg(ph->p_pa, ph->p_memsz, ph->p_offset); // call the entry point from the ELF header // note: does not return! ((void (*)(void)) (ELFHDR->e_entry))();

Exercise 4-6

- Ex 4: Understand why pointer.c works like that and read about ELF header and the ELF file..
 - https://en.wikipedia.org/wiki/Executable_and_Linkable_Format
- Ex 5: Use si, ni to follow the instructions after changing 0x7c00 to others, e.g., 0x6b00 or something else..

\$(OBJDIR)/boot/boot: \$(BOOT_OBJS) @echo + ld boot/boot \$(V)\$(LD) \$(LDFLAGS) -N -e start -Ttext 0x7C00 -o \$@.out \$^ \$(V)\$(OBJDUMP) -S \$@.out >\$@.asm \$(V)\$(OBJCOPY) -S -0 binary -j .text \$@.out \$@ \$(V)perl boot/sign.pl \$(OBJDIR)/boot/boot

• Ex 6: practice gdb commands

GDB Command for Reading Memory

- x/100wx [address or register]
 - Examine
 - **100** values
 - sized as word (w, 4 bytes)
 - b byte
 - g-8 bytes
 - In hexadecimal (x)
 - d decimal

>>> x/10	00wx 0x7c00		
0x7c00:	0xc031fcfa	0xc08ed88e	0x64e4d08e
0x7c10:	0x64e6d1b0	0x02a864e4	0xdfb0fa75
0x7c20:	0x0f7c6416	0x8366c020	0x220f01c8
0x7c30:	0xb8660008	0xd88e0010	0xe08ec08e
0x7c40:	0x007c00bc	0x00c0e800	0xfeeb0000
0x7c50:	0x00000000	0x0000ffff	0x00cf9a00
0x7c60:	0x00cf9300	0x7c4c0017	0xba550000
0x7c70:	0x83ece589	0x403cc0e0	0xc35df875
0x7c80:	0x0c5d8b53	Øxffffe1e8	0x01f2baff
0x7c90:	0xc3b60fee	0x0feef3b2	0xf4b2c7b6
0x7ca0:	0x10e8c1f5	0xeec0b60f	0xb218ebc1
0x7cb0:	0xb0eee0c8	0xeef7b220	0xffffade8
0x7cc0:	0x000080b9	0x01f0ba00	0xf2fc0000
0x7cd0:	0xe58955c3	0x0c7d8b57	0x10758b56
0x7ce0:	0x0109eec1	0xe38146df	0xfffffe00
0x7cf0:	0x81534656	0x000200c3	0xff7ee800
0x7d00:	0x658deaeb	0x5f5e5bf4	0x8955c35d
0x7d10:	0x10006800	0x00680000	0xe8000100
0x7d20:	0x810cc483	0x0100003d	0x4c457f00
0x7d30:	0x0001001c	0x0000988d	0xb70f0001
0x7d40:	0x05e0c100	0x3903348d	0xff1673f3
0x7d50:	0xf473ff20	0xe8ec73ff	0xffffff75
0x7d60:	0x1815ffe6	0xba000100	0x00008a00
0x7d70:	0xb8ef66ff	0xffff8e00	0xfeebef66
0x7d80:	0x00000000	0×00000000	0x00000000

0xfa7502a8 0x010f60e6 0x7c32eac0 0xd08ee88e 0x00000000 0x0000ffff 0x000001f7 0x57e58955 0x01b00000 0xb2d889ee 0x83d888f6 0x087d8bff 0x5d5f5b6d 0x085d8b53 0x1273fb39 0x5a58ffff 0x6a5356e5 0xfffffb1 0xa1387546 0x01002c05 0xc3830473 0xeb0cc483 0xff8a00b8 0x00000000 0x00000000

Coding Convention (CODING)

- No space after a function name in a call
 - cprintf("asdf") GOOD
 - cprintf ("asdf") NO
- One space after if/for/while/switch
 - if (a == 1) { GOOD
 - if(a==1) { NO
- function_and_variable_names_look_like_this
 - NoCamelCase
- Macros are ALL UPPERCASE
 - e.g., SEG()

Coding Convention (CODING)

- Pointer types includes a space before *
 - (uint32_t *) GOOD
 - (uint32_t*) NO
- Use '//' for your comment
 - All imported comments are /**/, so we can distinguish yours from those
 - FYI, Linux Kernel uses /**/...
- Function with no args
 - f(void), not f();

Coding Convention (CODING)

• Function definition

- Insert newline between the return type and function name
- This will make finding function definition easy
- E.g., find the definition of mon_kerninfo would be:
- '^mon_kerninfo' in regexp.

int
mon_kerninfo(int argc, char **argv, struct Trapframe *tf)

os2 ~/cs444/s21/os2-lab1-Rogersyp 106% grep -nr ^mon_kerninfo
kern/monitor.c:43:mon_kerninfo(int argc, char **argv, struct Trapframe *tf)
os2 ~/cs444/s21/os2-lab1-Rogersyp 107%