

# CS 162

# Intro to Computer Science II

Lecture 24

Linked List

Assignment 4 Help

Final remarks

3/15/24



**Oregon State**  
University

# Odds and Ends

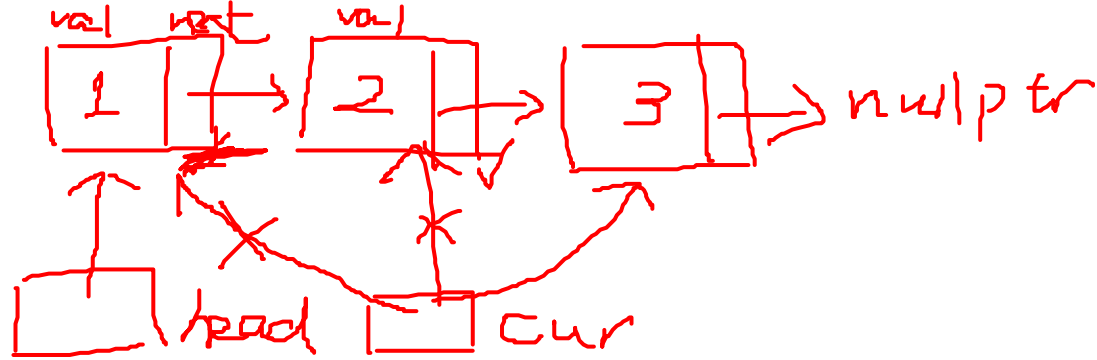
- Due reminder:
  - Quiz 5 due Sunday midnight via Canvas – open after today's lecture
  - Assignment 4 due Sunday midnight via TEACH
    - Grace days are allowed
- Today is the last day to demo:
  - Assignment 3 without late demo penalty
  - Assignment 1&2 with 30% late demo penalty
- Final Exam:
  - Wednesday 3/20 at 12pm at LINC 200

# Today's topic(s)

- Linked list
- Final exam review

Size 3

## In class activity



- Use the code provided on Canvas, complete the following tasks:
  - Task 1: What does the code do? (Hint: Trace through the code by drawing the picture out)
  - Task 2: Write code to print the list you just created. Trace the code you wrote to verify
    - Hint: Use while loop and Node\* current
  - Task 3: Delete the list you just created. Trace the code you wrote to verify
    - Hint: You might need another Node\*

# Pros and Cons of Singly Linked List

- Pros
  - Easy to implement
  - Insertion and deletion of elements can be done easily and doesn't requires movement of all elements compared to an array
  - Can allocate or deallocate memory easily during its execution
- Cons
  - Uses more memory when compared to an array
  - No random access
  - Traversing in reverse is not possible for singly linked list

# Today's topic(s)

- Linked list
- Final exam review

# Final Exam

- Weight: 15% of course grade
- Time: Wednesday 12:00 – 12:50 pm
- Where: LINC 200
- Close book, close notes, no calculator
- Scratch paper will be provided if needed
- Bring pen/pencil, and your photo ID
- Question types:
  - T/F, multiple choice
  - Similar as the midterm exam 😊
- Question amount: ~40

# Coverage

- Non-cumulative
- Emphasis on material covered after Midterm (90%)
  - Lecture 14-24 (start from shallow vs. deep copy)
  - Lab 6 – 10
  - Worksheet 6-10
  - Assigned Reading
  - Assignment 3-4
- General coverage of earlier topics (10%)



# Topics

- Shallow vs. Deep copy
- Big 3 and their usage
- Inheritance
- Upcasting vs. downcasting
- Polymorphism
- Virtual vs. pure virtual
- Abstract class
- Function/class templates
- Standard Template Class (STL)
  - vector
- Containers
- Linked List (singly)
- Recursion

# Study Guide

- Take the practice exam and time yourself
- Lecture slides 14-24
- Quiz 3-5
- Worksheet 6-10
- Lab 6-10
- Assignment 3-5
- Assigned readings

# Winter 2023 Exam Review

# You have learned MANY things from CS 162

- Pointers
- Memory model
  - Stack vs. heap
- Dynamic Arrays
- Structs
- File separation
  - .h .cpp
  - Header guards
- Makefile
  - Compilation process
- File I/O
- Object Oriented Programming
  - Encapsulation
- Struct vs. Class
- C++ Classes
  - Access specifiers: private, public, protected
  - Accessor and Mutator functions
  - this keyword
  - Constructors: default vs. non-default
  - const
  - Big Three
    - Copy constructor
    - Assignment operator overload
    - Destructor
  - Class composition vs. class inheritance
  - Polymorphism

# You have learned MANY things from CS 162

- Template
  - STL
- Containers
  - Linked list vs. array
  - Vector
- Recursion

# Be Confident...



Now you are able to...

- Design and implement programs that require:
  - multiple classes and structures
  - hierarchies of classes that use inheritance and polymorphism
  - an understanding of abstraction, modularity and separation of concerns
- Construct and use basic linear structures (arrays, stacks, queues, and various linked lists) in programs, and be able to describe instances appropriate for their use.
- Develop test-data sets and testing plans for programming projects.
- Produce recursive algorithms, and choose appropriately between iterative and recursive algorithms.

# Final Remarks...

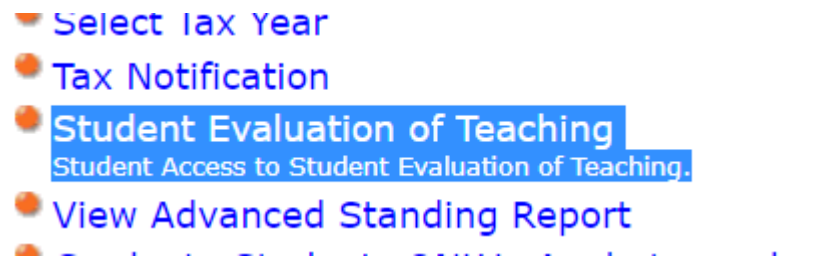
- Thank you so much for your commitment to this course

- What's next ?

- CS 261: Data Structure
- ECE/CS 271: Computer Architecture and Assembly Language
- CS 290: Web Development

- Future improvements?

- Canvas SLE →



- ULA position

- Contact me! And apply through: <https://jobs.oregonstate.edu/postings/140560>

# Final Remarks...

- Submit all your work by the deadline
  - Assignment 4, Quiz 5
- Take the Final Exam on Wednesday
  - Bring your photo ID
- Grade disputation:
  - By 3/23 6pm



# Assignment 4 Q&A

# \*Additional topic(s)

- Complexity Analysis

\*Note: this will not be in the final

# How to compare/describe algorithms

- We have different data structures and sorting algorithms, how to compare them?
- We want a way to characterize runtime or memory usage that is completely **platform-independent**
  - i.e. does not depend on hardware, operating system, programming language, etc.

# Complexity Analysis

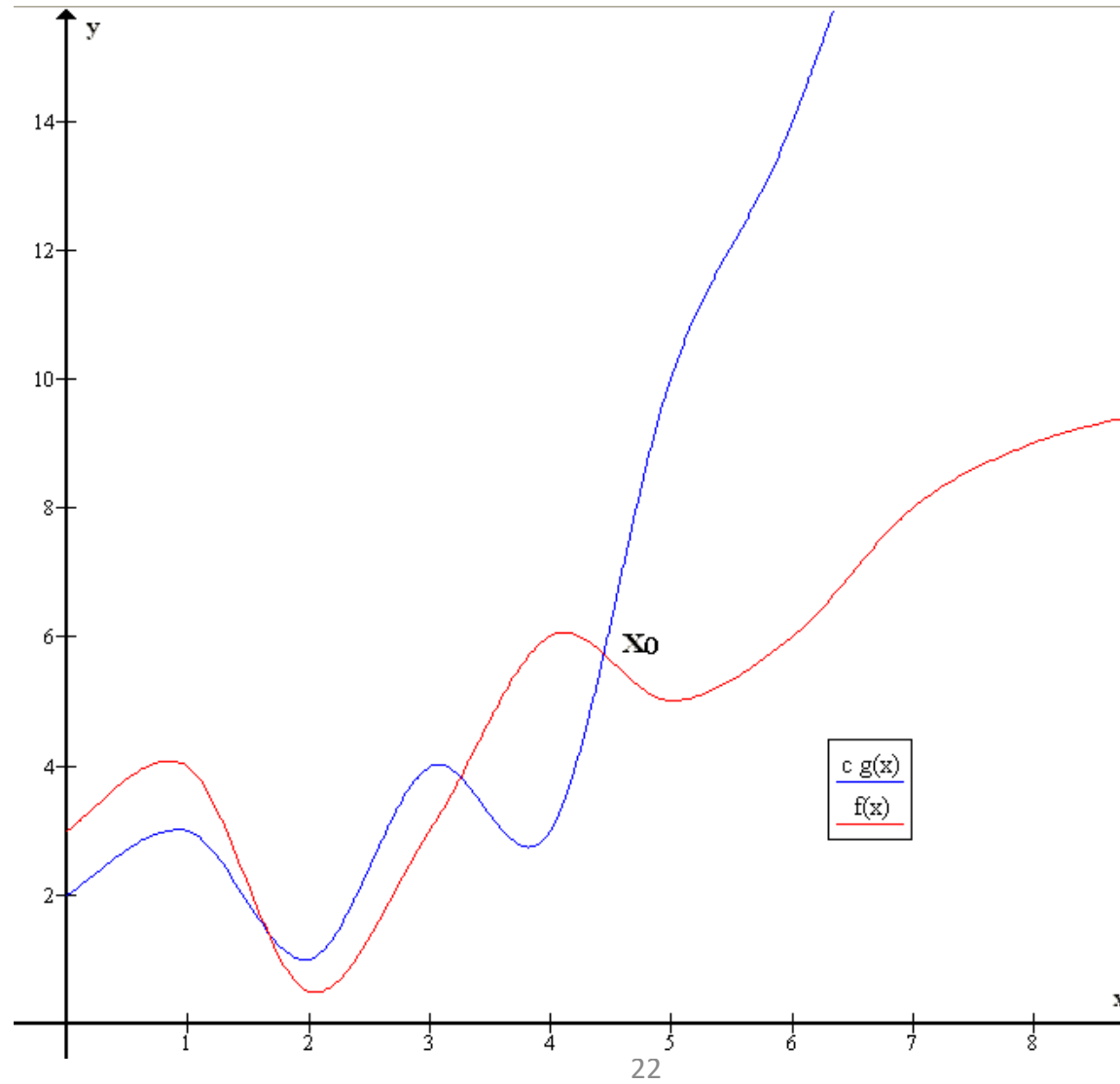
- Use **Complexity Analysis** to help make platform-independent comparisons of data structures
  - Refer to as **Big O**
- Allow us to assess a data structure or algorithm's resource usage (i.e., runtime and memory consumption) in an abstract way
- To do this, we describe how a data structure's or algorithm's runtime or memory usage changes relative to a change in the input size (**n**)

# Big O

- We use **Big O notation** to assess a data structure or algorithm's performance.
- Big O notation: a tool for characterizing a function in terms of its **growth rate**
  - Indicate an **upper bound** on the function's growth rate, known as **growth order**

# Big O

$g(x)$  provides an upper bound on  $f(x)$

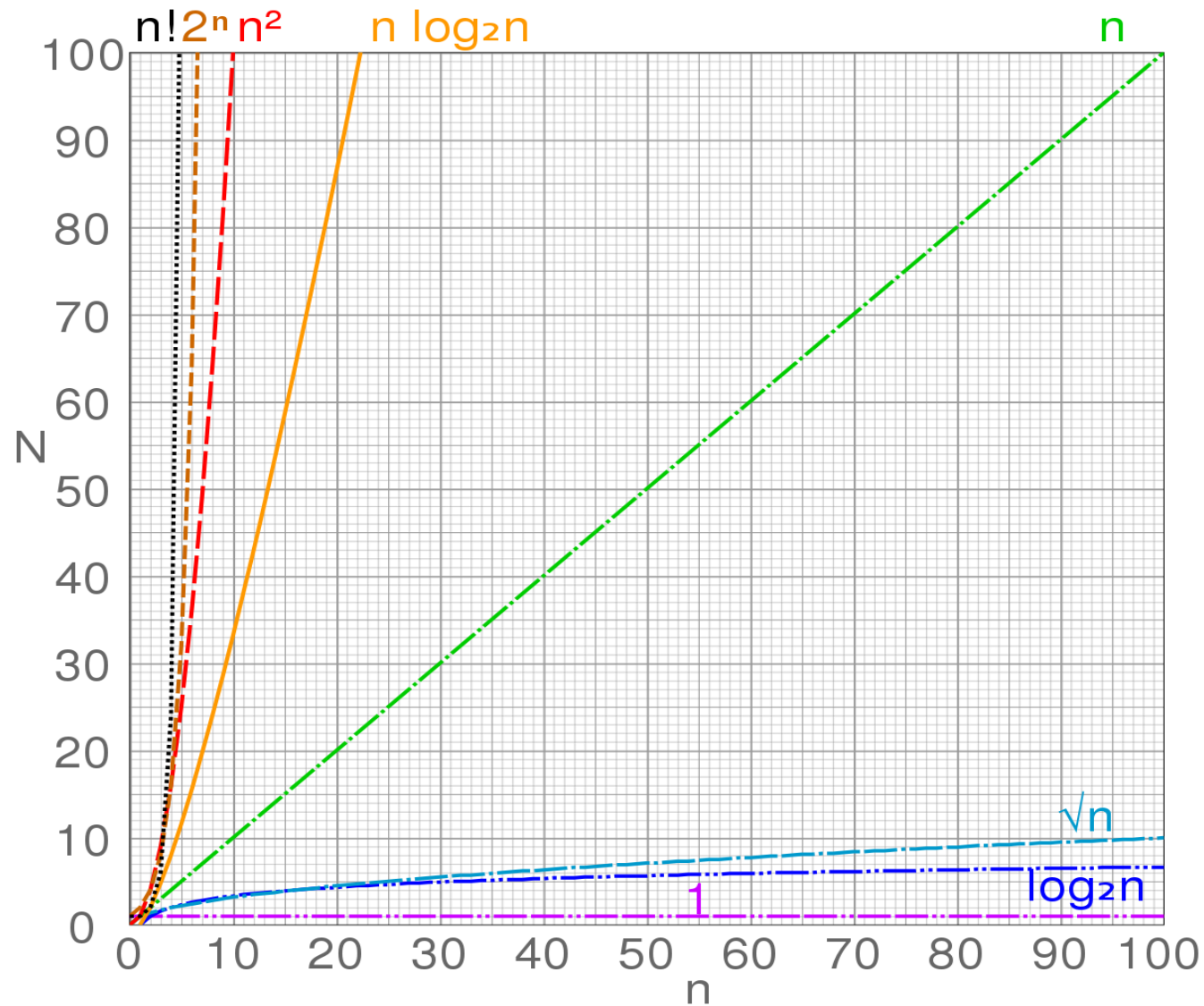


$g(x)$  is  $O(f(x))$

# Big O

- To assess a data structure or algorithm's complexity, we will compute a growth order for its runtime (or memory usage) as a function of the input size  $n$
- Importantly, we want to describe how data structures behave **in the limit, as  $n$  approaches  $\infty$  (infinity)**

# Common growth order functions





# Common growth order functions

- $O(1)$  – constant complexity
- $O(\log n)$  – log-n complexity
- $O(\sqrt{n})$  – root-n complexity
- $O(n)$  – linear complexity
- $O(n \log n)$  – n-log-n complexity
- $O(n^2)$  – quadratic complexity
- $O(n^3)$  – cubic complexity
- $O(2^n)$  – exponential complexity
- $O(n!)$  – factorial complexity

# Big O

- Consider this example...

```
int sum = 0;
for (i = 0; i < n; i++) {
    sum += array[i];
}
return sum;
```

- This function is summing an array of n integers
- What's the run-time complexity of the function?

# Big O example

```
int sum = 0;
for (i = 0; i < n; i++) {
    sum += array[i];
}
return sum;
```

- The instruction `int sum = 0;` executes in some constant time  $c_1$  independent of  $n$
- Each iteration of the loop executes in some constant time  $c_2$ , and this happens  $n$  times
- The return statement executes in some constant time  $c_3$  independent of  $n$
- So runtime is  $c_1 + c_2 * n + c_3$
- $c_1$ ,  $c_2$ , and  $c_3$  depend on the particular computer running this function, so we ignore them to figure out run-time complexity
- Thus, this function grows on the order of  $n$ , a.k.a. its run-time complexity is  **$O(n)$**

# Determining a program's complexity

```
node* push (node * head, int val) {  
    node *temp = new node;  
    temp->val =val;  
    temp->next = head;  
    head = temp;  
    return head;  
}
```

- Every instruction in this function executes in some constant time, independent of n
- Thus we ignore them to figure out runtime complexity.
- Complexity:  $O(c_1+c_2+c_3+c_4+c_5) = \mathbf{O(1)}$

# Dominant components

- When a growth order function has additive terms, one of those will dominate the others
  - Specifically, function  $f(n)$  dominates  $g(n)$  if  $n_0: n > n_0, f(n) > g(n)$
- In these cases, we simply ignore the non-dominant terms
  - i.e.  $n^2 - n, n^2$  dominates  $n$ , so we ignore  $n$ , and we say this complexity is  $O(n^2)$

# More examples

- Loops are one of the main determinants of a program's complexity

- ```
for (int i = 0; i < n; i++) {  
    ...  
}
```

- ```
for (int i = n; i > 0; i/=2) {  
    ...  
}
```

- ```
for (int i = 0; i*i < n; i++) {  
    ...  
}
```

# More examples

- ```
for (int i = 0; i < n; i++) {  
    for (int j = 0; j < n; j++) {  
        ...  
    }  
}
```

- ```
for (int i = n; i > 0; i/=2) {  
    for (int j = 0; j < n; j++) {  
        ...  
    }  
}
```

# Real-world Consideration

- Your program will only perform as well as your design
  - Constant factors can still play a part
- Suppose you have two algorithms...
  - Algorithm A)  $1,000,000n \rightarrow O(n)$
  - Algorithm B)  $2n^2 \rightarrow O(n^2)$
  - Which one is better?
    - It depends